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Verifying enclave systems with Serval

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Enclave monitors are hard to get right

- Correctness of enclave monitor code is critical for security
- Many different kinds of bugs are security vulnerabilities:
 - Low-level bugs: e.g., buffer overflow or division-by-zero
 - Logic bugs: implementation does something unintended
 - **Design bugs**: intended design of the system is not secure
- Each can be exploited to compromise the entire system

Eliminating bugs with formal verification

• Goal: prove absence of low-level, logic, and design bugs

- **Approach**: Use automated verification techniques
 - Low proof burden: symbolic evaluation / SMT solvers
 - Bounded loops in code, likely true of monitors
 - Limitations: no concurrency or side channels

Challenges

- Difficulty of building verifiers
 - Need detailed RISC-V machine model
 - Need to reason at ISA level
- Difficulty of scaling to practical systems
 - Symbolic evaluation
 - SMT solving

Serval: A framework for verifying low-level systems

- Built on top of Rosette
- Lift ISA interpreter into verifier
 - Easier to write and test
 - Supports LLVM IR and RISC-V
- Use symbolic profiling to identify verification bottlenecks
- Use symbolic optimizations to scale verification

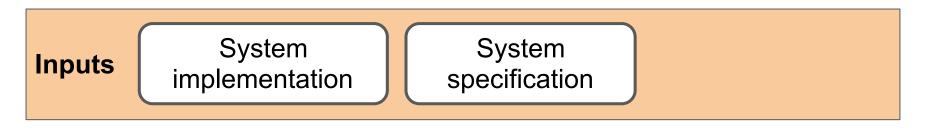
Main Results

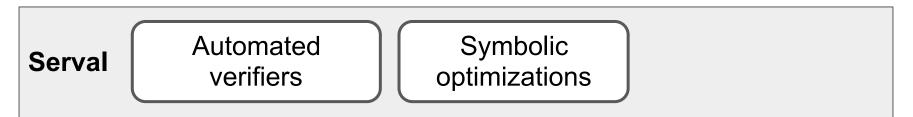
- Applied to CertiKOS (PLDI'16) and Komodo (SOSP'17), previously manually verified using Coq and Dafny
- Found and fixed 15 Linux BPF JIT bugs, all now upstreamed.
 - <u>https://git.kernel.org/linus/1e692f09e091</u>
 - https://git.kernel.org/linus/46dd3d7d287b
 - https://git.kernel.org/linus/68a8357ec15b
 - https://git.kernel.org/linus/6fa632e719ee
- Work in progress: identified implementation and design bugs in Keystone

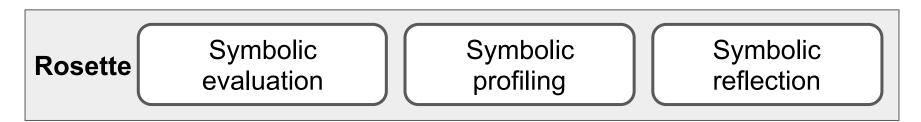
Outline

- Verification stack
- Workflow
- Demo

Verification stack

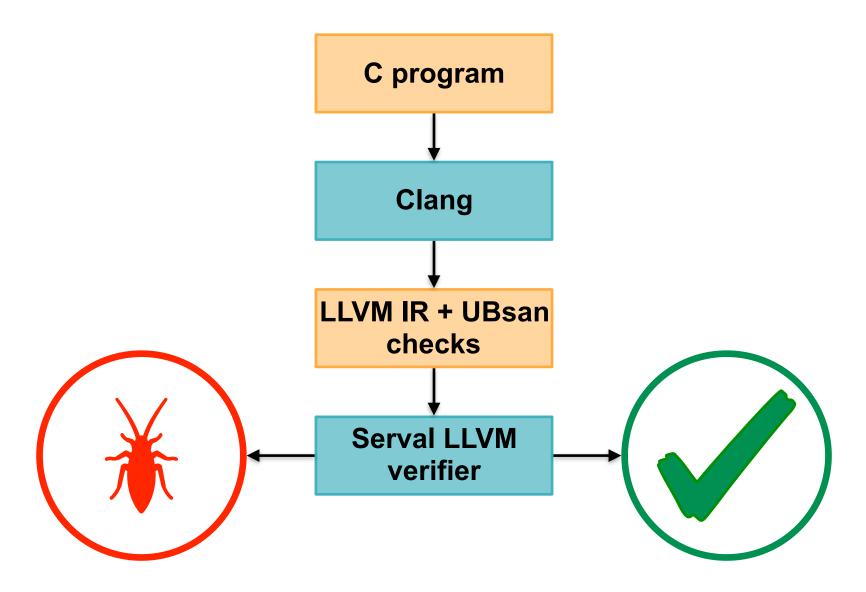




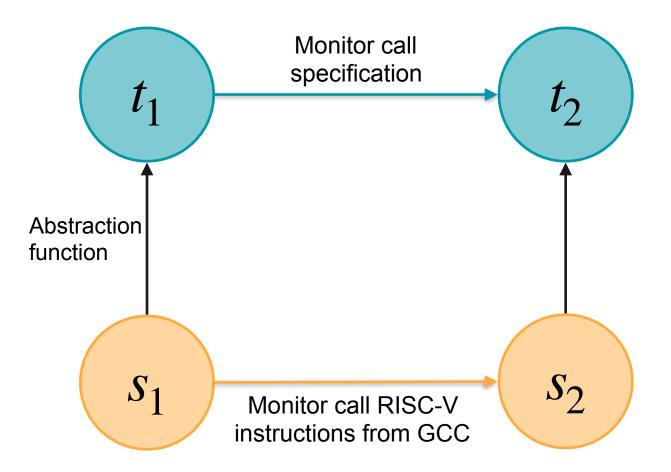




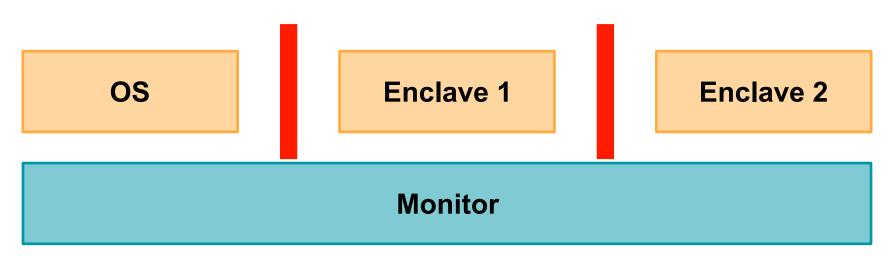
[1/3] Proving absence of low-level bugs



[2/3] Proving functional correctness



[3/3] Proving noninterference



- Example: bogus monitor call that returns enclave secrets
- Integrity OS should not be able to modify enclavevisible state
- Confidentiality Behavior of OS is independent of enclave secrets

Demo: Komodo

- A verified software enclave monitor
- We have ported to RISC-V and verified using Serval

- Demonstration:
 - Low-level buffer overflow vulnerability

Demo

Conclusion

- Automated verification is effective at eliminating bugs in low-level systems
- If you are building enclave systems, talk to us!
- Paper to appear at SOSP'19
- Code will be released shortly

https://serval.unsat.systems/ UNSAT